

## 19 Undirected graphs

A *Graph* is a pair  $(V, E)$  of *vertices* and *undirected edges*.

There is a notation which we don't use much:

$$V^{(2)} = \{\{u, v\} : u, v \in V \wedge u \neq v\}$$

the set of all possible undirected edges with vertices in  $V$ .

So:

**(19.1) Definition** *An (undirected) graph is a pair  $(V, E)$  of vertices and (undirected) edges where  $V$  is any finite set and  $E \subseteq V^{(2)}$ .*

In fact, we shall represent graphs as ‘bidirected graphs,’ which are identical to directed graphs, except that whenever there is an edge  $e = (u, v)$ , there is another edge  $e' = (v, u)$ , and every edge carries an extra ‘inverse’ component;  $e'$  is the inverse of  $e$  and  $e$  is the inverse of  $e'$ .

For the purposes of this module all graphs will be input as bidirected graphs, as, for example

```
6 14
0  3  4  2  5
1  2  3  5
2  2  4  0
3  2  1  5
4  2  0  2
5  3  3  1  0
```

After a graph has been read in, it is necessary to install the ‘inverse’ links in the edges.

The code samples include a program `bcc.c` whose purpose is to calculate the biconnected components of  $G$ , see below. It has a routine `link_inverse_edges` which installs these links:

```
void link_inverse_edges ( GRAPH * graph )
{
    int i;
    for (i=0; i<graph->n; ++i)
    {
        int j;
        EDGE * e = graph->an_edge[i];
        for (j=0; j<graph->out_deg[i]; ++j)
        {
            if ( e->inverse == NULL )
            {
                int k;
                EDGE * ee = graph->an_edge[e->to];
                int found = 0;
                for ( k=0; k<graph->out_deg[e->to] && (!found); ++k )
                {
                    if ( ee->to == i )
```

```

        { e->inverse = ee; ee->inverse = e; found = 1; }
        ee = ee->next;
    }
}
e = e->next;
}
}
}

```

This code is ‘inefficient’ because in order to set up the inverse link from an edge, it inspects all the edges in the graph. In other words, the runtime is  $\Omega(m^2)$  where  $m$  is the number of edges.

It is possible to accomplish this in linear time,  $O(m)$ , using a variant of lexical sort, and the third programming assignment is to replace the above piece of code by a more efficient one (of course any improvement in the runtime will not be noticeable).