

16. Irrep(S_n) = { Specht reps V_λ | $\lambda \vdash n$ }

Our today's goal is to prove the statement from the title, which was already illustrated for small n . The proof is decomposed into a list of lemmas.

Recall that a partition $\lambda = (\lambda_1, \dots, \lambda_k)$ of n , written $\lambda \vdash n$, is a decomposition $n = \lambda_1 + \dots + \lambda_k$, $\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_k > 0$, $\lambda_i \in \mathbb{N}$.
Convention: $\lambda_i = 0$ for $i > k$.

Def: $\lambda \supseteq \lambda'$, for $\lambda, \lambda' \vdash n$, means $\lambda_1 \geq \lambda'_1$
 $\uparrow$
dominates
 $\lambda \supset \lambda'$
$$\begin{cases} \lambda_1 + \lambda_2 \geq \lambda'_1 + \lambda'_2 \\ \lambda_1 + \lambda_2 + \lambda_3 \geq \lambda'_1 + \lambda'_2 + \lambda'_3 \\ \dots \end{cases}$$

Ex: $(3, 1, 1) \supseteq (2, 2, 1)$

$(4, 1, 1)$ and $(3, 3)$: no domination.

Rmk: Relation $\supseteq$ is a partial order.

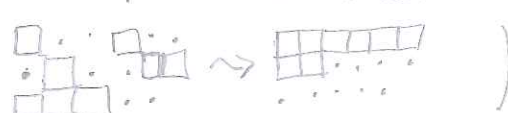
Lemma 20: $\lambda, \lambda' \vdash n$, $T \in Y_{T_\lambda}$, $T' \in Y_{T_{\lambda'}}$.

(a) If all i & j from the same row of T' lie in different columns of T , then $\lambda \supseteq \lambda'$.

(b) If moreover $\lambda = \lambda'$, then $T' \sim T'' = \sigma \cdot T$ for some $\sigma \in C_T$.

□ (a) Put $X_i = \{j \mid j \text{ lies in the first } i \text{ rows of } T'\}$. $\#X_i = \lambda'_1 + \dots + \lambda'_i$.

A column of T contains $\leq i$ numbers from X_i . Move these numbers to the first i rows of T by column permutations.

(Think of the tetris game: )

But in these i rows there must be $\lambda_1 + \dots + \lambda_i$ numbers. Thus $\lambda_1 + \dots + \lambda_i \geq \lambda'_1 + \dots + \lambda'_i$.

(b) Further column permutations of T move all letters from row $p \leq i$ of T' to rows $\leq p$ of T . Since $\lambda_p = \lambda'_p$ for all p , in fact the row p of T' will land precisely inside the row p of T . For $i = \# \text{Rows}(\lambda)$, this yields the desired $T'' \in Y_{T_\lambda}$.

□

Lemma 2.1: $\lambda, \lambda' \vdash n$, $T \in Y_{T_\lambda}$, $T' \in Y_{T_{\lambda'}}$.

(a) If $\sum_{\sigma \in G_T} \text{sgn}(\sigma) \sigma \cdot e_{[T']} \neq 0$, then $\lambda \triangleright \lambda'$.

recall the def. of E_T !

(b) If $\lambda = \lambda'$, then $\sum_{\sigma \in G_T} \text{sgn}(\sigma) \sigma \cdot e_{[T']} = \pm E_T$ or 0 .

□ (a) To deduce this from L. 2.0, one needs to show that no i, j from the same row of T' lie in the same column of T . Such i, j would satisfy $(ij) \cdot [T'] = [T']$ & $(ij) \in G_T$.

But then $C_T = \{\tau_1, \dots, \tau_k, \tau_k(ij), \dots, \tau_k(ij)\}$ for some $\tau_k \in S_n$

(coset decomposition).

$$\text{so, } \sum_{\sigma \in G_T} \text{sgn}(\sigma) \sigma \cdot e_{[T']} = \sum_{t=1}^k \left(\text{sgn}(\tau_t) e_{\tau_t \cdot [T']} + \underbrace{\text{sgn}(\tau_t(ij))}_{=-\text{sgn}(\tau_t)} \underbrace{e_{\tau_t(ij) \cdot [T']}}_{=\tau_t \cdot [T']} \right) = 0,$$

contradiction.

(b) L. 2.0 (b) allows one to continue the previous argument if $\lambda = \lambda'$: $\exists T' \sim T'' = (\sigma')^{-1} \cdot T \Rightarrow$

$$\sum_{\sigma \in G_T} \text{sgn}(\sigma) \sigma \cdot e_{[T']} = \sum_{\sigma'' \in G_{T''}} \text{sgn}(\sigma'' \sigma') \sigma'' \sigma' \cdot e_{[T'']} = \text{sgn}(\sigma') \sum_{\sigma'' \in G_{T''}} \text{sgn}(\sigma'') \sigma'' \cdot e_{[\sigma'^{-1} T']} = \text{sgn}(\sigma') E_T. \quad \square$$

Since the $e_{[T']}, T' \in Y_{T_\lambda}$, span $\mathbb{C} M_\lambda$, one gets

Cor 2.2: $\lambda \vdash n$, $T \in Y_{T_\lambda}$, $v \in \mathbb{C} M_\lambda \Rightarrow \sum_{\sigma \in G_T} \text{sgn}(\sigma) \sigma \cdot v = \lambda E_T$ for some $\lambda \in \mathbb{C}$.

Lemma 2.3: $\lambda, \lambda' \vdash n$, $\Phi \in \text{Hom}_{S_n}(\mathbb{C} M_\lambda, \mathbb{C} M_{\lambda'})$, Φ non-zero on $V_\lambda \subseteq \mathbb{C} M_\lambda$.

(a) Then $\lambda \triangleright \lambda'$.

(b) If moreover $\lambda = \lambda'$, then Φ restricts to $\lambda \text{Id}_{V_\lambda}$, for some $\lambda \in \mathbb{C}$.

This statement hopefully recalls you Schur's lemma!

□ (a) For $T \in Y_{T_\lambda}$, $E_T \notin \ker \Phi$: otherwise $V_\lambda = \sum_{\sigma \in S_n} \mathbb{C} \sigma \cdot E_T \subseteq \ker \Phi$, since $\ker \Phi$ is a subrep. But Φ is non-zero on V_λ .

$$\text{So } 0 \neq \Phi(E_T) = \Phi\left(\sum_{\sigma \in G_T} \text{sgn}(\sigma) \sigma \cdot e_{[T]}\right) \stackrel{\Phi \text{ is } \mathbb{C} \text{-} S_n \text{-linear}}{=} \sum_{\sigma \in G_T} \text{sgn}(\sigma) \sigma \cdot \Phi(e_{[T]})$$

$$\Rightarrow 0 \neq \sum_{\sigma \in G_T} \text{sgn}(\sigma) \sigma \cdot e_{[T']} \text{ for some } T' \in Y_{T_{\lambda'}}$$

$$\Rightarrow \lambda \triangleright \lambda', \text{ by L. 2.2.}$$

(b) By Cor 2.2, $\Phi(E_T) = \lambda E_T$ for some $\lambda \in \mathbb{C}$. Then $\Phi(\sigma \cdot E_T) = \sigma \cdot \Phi(E_T) = \lambda \sigma \cdot E_T$ for all $\sigma \in S_n$, thus Φ restricts to λId on $V_\lambda = \sum_{\sigma \in S_n} \mathbb{C} \sigma \cdot E_T$. $\square$

Thm 10: (1) $\forall \lambda \vdash n$, V_λ is irreducible

(2) $V_\lambda \cong V_{\lambda'} \Rightarrow \lambda = \lambda'$

(3) $\text{Irrep}(S_n) = \{V_\lambda \mid \lambda \vdash n\}$

(4) $\mathbb{C}M_\lambda \cong V_\lambda \oplus (\bigoplus_{\lambda' \triangleright \lambda} m_{\lambda, \lambda'} V_{\lambda'})$ for some $m_{\lambda, \lambda'} \in \mathbb{N} \cup \{0\}$.

□ (1) Take a subrep. W of $V_\lambda \subseteq \mathbb{C}M_\lambda$, $W \neq \{0\}$.

By Maschke's thm, $\mathbb{C}M_\lambda = W \oplus U$, as S_n -reps.

Consider $\varphi \in \text{Hom}_{S_n}(\mathbb{C}M_\lambda, \mathbb{C}M_\lambda)$, given by $\varphi(w, u) = (w, 0)$.

By L. 23.19, φ restricts to Id_W . But $\varphi(v) = v \Leftrightarrow v \in W$, so $V_\lambda = W$.

(2) $\exists \varphi \in \text{Iso}_{S_n}(V_\lambda, V_{\lambda'}) \Rightarrow$ it extends to $\hat{\varphi} \in \text{Hom}_{S_n}(\mathbb{C}M_\lambda, \mathbb{C}M_{\lambda'})$, as in (1) $\Rightarrow \lambda \triangleright \lambda'$. Starting from φ^{-1} , one gets $\lambda' \triangleright \lambda$.
L. 23.1a) so $\lambda = \lambda'$.

(3) $\Leftarrow$ (1) & (2) & $\# \text{Irrep}(S_n) = \# \text{Conj}(S_n) \stackrel{\text{Thm 7}}{=} \# \{\lambda \vdash n\}$.

(4) By (3), $\mathbb{C}M_\lambda \cong \bigoplus_{\lambda' \triangleright \lambda} m_{\lambda, \lambda'} V_{\lambda'}$. Now, complete inclusions $V_{\lambda'} \rightarrow \mathbb{C}M_\lambda$ to S_n -morphisms $\mathbb{C}M_{\lambda'} \rightarrow \mathbb{C}M_\lambda$ as in (1), and use L. 23.1a). □

17. An algorithm for constructing the character table of S_n for $\forall n$.
(4) allows one to construct inductively irreps of S_n starting with V_n , using Prop. 24. $\uparrow$ w.r.t. $\triangleright$

Ex.: $6 \triangleright 5, 1 \triangleright 4, 2 \triangleright 3^2 \triangleright 3, 2, 1 \triangleright 2^3 \triangleright 2^2, 1^2 \triangleright 2, 1^4 \triangleright 1^6$

Prop. 24: $\forall \lambda = (\lambda_1, \dots, \lambda_k) \vdash n$, $\lambda' = (\lambda'_1, \dots, \lambda'_{k'}) \vdash n$, $\sigma \in S_n$ of cycle type λ' ,

$\chi^{\mathbb{C}M_\lambda}(\sigma) = \text{coefficient of } x_1^{\lambda_1} \dots x_k^{\lambda_k} \text{ in } \prod_{j=1}^{k'} (x_1^{\lambda'_j} + \dots + x_k^{\lambda'_j})$

Ex.: $n=4$, $\lambda' = (3, 1)$, $(x_1^3 + x_2^3)(x_1 + x_2) = x_1^4 + x_2^4 + x_1^3 x_2 + x_2 x_1^3 \Rightarrow$

$\chi^{\mathbb{C}M_\lambda}((123)) = \begin{cases} 1, & \text{if } \lambda = (4) \text{ or } (3, 1), \\ 0, & \text{otherwise,} \end{cases}$

$k=2$ suffices
 $\uparrow$
Cre 25

Rmk: $\prod_i x_i^{\lambda_i}$ is a symmetric polynomial in the x_i .

In fact, the connection between S_n -reps & symmetric polynomials is very deep!

Proof of Prop. 24: $\mathbb{C}M_\lambda$ is a permutation rep. of S_n .

! It is in general different from the permutation rep.

$$V^{\text{perm}} = \bigoplus_{i=1}^n \mathbb{C}e_i, \sigma \cdot e_i = e_{\sigma(i)}.$$

Then $\chi^{\mathbb{C}M_\lambda}(\sigma) = \#(M_\lambda)^\sigma = \# \{ [T] \in M_\lambda \mid \sigma \cdot [T] = [T] \}$
↑ of cycle type λ'

$$= \# \{ T \in Y_{T_\lambda} \mid \forall i, i \text{ \& } \sigma(i) \text{ are in the same row of } T \} \quad \text{row reordering}$$

$$= \# \{ \varphi: \{1, \dots, k'\} \rightarrow \{1, \dots, k\} \mid \forall i, \lambda_i = \sum_{\varphi(j)=i} \lambda'_j \} \quad (*)$$

$$= \# \{ \varphi: \{1, \dots, k'\} \rightarrow \{1, \dots, k\} \mid \prod_i x_i^{\lambda_i} = \prod_j x_{\varphi(j)}^{\lambda'_j} \}$$

here $\varphi(j)=i \Leftrightarrow j^{\text{th}}$ cycle of σ lies inside i^{th} row of T

$$= \sum_{\varphi} (\text{coef. of } \prod_i x_i^{\lambda_i} \text{ in } \prod_j x_{\varphi(j)}^{\lambda'_j})$$

seen as a polynomial

$$= \text{coef. of } \prod_i x_i^{\lambda_i} \text{ in } \sum_{\varphi} \prod_j x_{\varphi(j)}^{\lambda'_j}$$

$$= \text{coef. of } \prod_i x_i^{\lambda_i} \text{ in } \prod_j (x_1^{\lambda'_j} + \dots + x_k^{\lambda'_j}) \quad \square$$

Cor 25: $\chi^{\mathbb{C}M_\lambda}(\lambda') \neq 0 \Rightarrow \lambda \triangleright \lambda' \Rightarrow k \leq k'$.

$\square \exists \varphi$ satisfying $(*) \Rightarrow \lambda'_1 \leq \lambda_{\varphi(1)} \leq \lambda_1,$

$$\lambda'_1 + \lambda'_2 \leq \underbrace{\lambda_{\varphi(1)} + \lambda_{\varphi(2)}}_{\text{if } \varphi(1) \neq \varphi(2)} \text{ or } \underbrace{\lambda_{\varphi(1)}}_{\text{if } \varphi(1)=\varphi(2)} \leq \lambda_1 + \lambda_2,$$

So one gets a triangular table.

Ex.: $n=4$

$\lambda \backslash \sigma$	Id	(12)	$(12)(34)$	(123)	(1234)
$\mathbb{C}M_4$	1	1	1	1	1
$\mathbb{C}M_{3,1}$	4	2	0	1	
$\mathbb{C}M_2$	6	2	2		
$\mathbb{C}M_{2,1^2}$	12	2			
$\mathbb{C}M_{1^4}$	24				

$\leq x_1^4$ always appears once

$$\mathbb{C}M_{2,1^2} \cong V^{\text{st}} \oplus V^{\text{reg}} \oplus W \oplus 2V^{\text{st}} \oplus V^{\text{tr}} \quad (*)$$

$\Leftarrow \# S_n$

$(x_1 + \dots + x_4)^4 = \sum_{(d_1, d_2, d_3, d_4)} \binom{4}{d_1, d_2, d_3, d_4} x_1^{d_1} \dots x_4^{d_4}$
 where $\binom{4}{d_1, d_2, d_3, d_4} = \frac{4!}{d_1! d_2! d_3! d_4!}$
 $(x_1^2 + x_2^2 + x_3^2)(x_1 + x_2 + x_3)^2 = x_1^4 + 2x_1^3 x_2 + 2x_1^2 x_2^2 + 2x_1^2 x_2 x_3 + \text{non-relevant terms}$
 $= x_1^4 + x_2^4 + 2x_1^2 x_2^2$

Our construction of character tables for S_3, S_4, S_5 relied on "happy coincidences": geometric actions, the irreducibility of certain "natural" reps, etc.

But now we have enough tools for designing an

Algorithm for constructing the char. table for any S_n

(1) List all partitions $\lambda \vdash n$.

(2) Order them: $\lambda^{(1)}, \lambda^{(2)}, \dots$ in such a way that

$$\lambda^{(j)} \triangleright \lambda^{(i)} \Rightarrow j \leq i$$

(e.g., take the opposite of the lexicographic ordering)

(3) Compute $\chi^{\chi^{\lambda^{(i)}}}$ for all $\lambda \vdash n$ using Prop. 24.

(4) For $i=1, 2, \dots$, in this order, compute

$$(a) m_{i,j} = (\chi^{\chi^{\lambda^{(i)}}}, \chi^{\chi^{\lambda^{(j)}}}) \quad \text{for all } j < i.$$

$$(b) \chi^{\chi^{\lambda^{(i)}}} = \chi^{\chi^{\lambda^{(i)}}} - \sum_{j < i} m_{i,j} \chi^{\chi^{\lambda^{(j)}}}.$$

The computation from the last point is motivated by Thm 10. (4).

This algorithm directly implies

Thm 11: The characters of all irreps of S_n are integer valued: $\chi^{\lambda}(\sigma) \in \mathbb{Z}$.

There are also methods for computing $\chi^{\lambda}(\sigma)$ directly:

A) **Frobenius formula:**

$$\chi^{\lambda}(\sigma) = \text{coefficient of } x_1^{\hat{\lambda}_1} \dots x_k^{\hat{\lambda}_k} \text{ in } \prod_{1 \leq i < j \leq k} (x_i - x_j) \cdot \prod_{j=1}^{k'} (x_1^{\lambda_j'} + \dots + x_k^{\lambda_j'}).$$

of cycle type λ'

$$\text{where } \hat{\lambda}_k = \lambda_k, \hat{\lambda}_{k-1} = \lambda_{k-1} + 1, \dots, \hat{\lambda}_1 = \lambda_1 + k - 1.$$

B) **Murnaghan - Nakayama rule**

in terms of skew Young tableaux.

We do not give proofs here.

Rmk: In contrast to what we have seen in the last several lectures, the reps of S_n over a finite field $\mathbb{F}_q$ are not completely understood.