

A bridge between knotted graphs and axiomatizations of groups

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OCAMI, Osaka

FMSP Lectures
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Based on:

✍ Qualgebras and Knotted 3-Valent Graphs, *to appear in Fundamenta Mathematica*

✍ (With S. Kamada) Alexander and Markov Theorems for Graph-Braids, *in progress*

A bridge between knotted graphs
and axiomatizations of groups



topology

A bridge between knotted graphs and axiomatizations of groups



algebra

topology

Qualgebras: A bridge between knotted graphs and axiomatizations of groups



algebra

topology

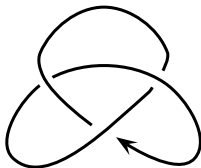
Part 1:

*How a Knot Theorist
Would Invent Qualgebras*

Knot diagrams: from illustration to manipulation

1926 K. Reidemeister:

$$\text{Knots} \cong \text{Diagrams} / \text{R-moves}$$

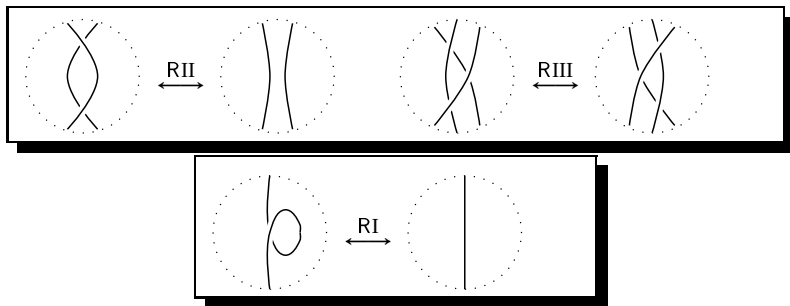


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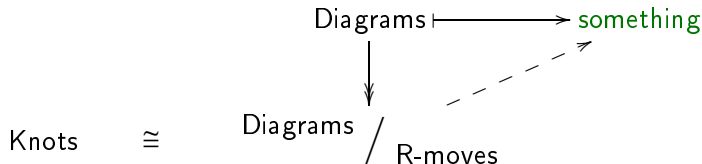


Knot diagrams: from illustration to manipulation

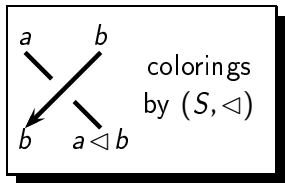
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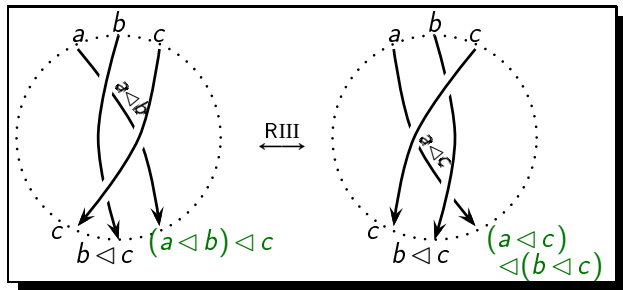
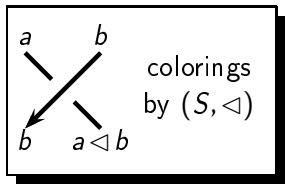
Combinatorial knot invariants:



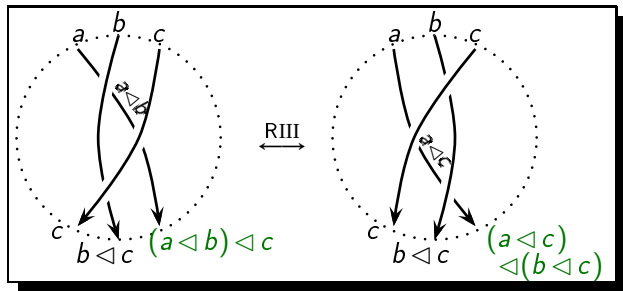
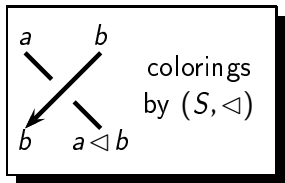
Quandles as an algebraization of knots



Quandles as an algebraization of knots

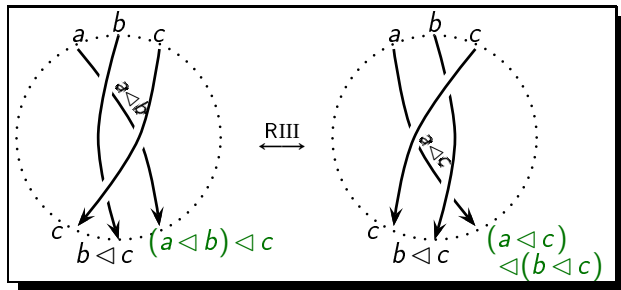
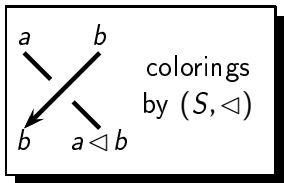


Quandles as an algebraization of knots



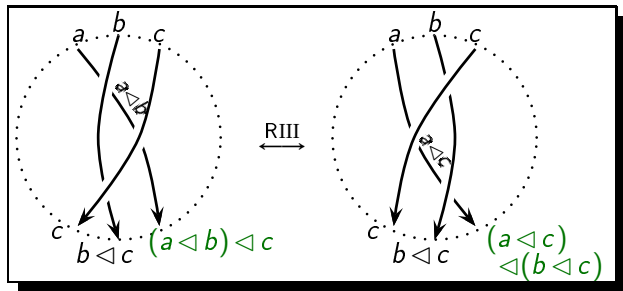
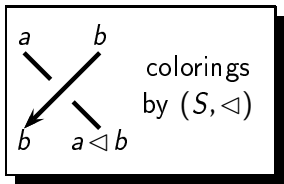
$$\text{RIII} \leftrightarrow (a \triangleleft b) \triangleleft c = (a \triangleleft c) \triangleleft (b \triangleleft c) \quad (\text{SD})$$

Quandles as an algebraization of knots



RIII	$\leftrightarrow$	$(a \triangleleft b) \triangleleft c = (a \triangleleft c) \triangleleft (b \triangleleft c)$	(SD)
RII	$\leftrightarrow$	$(a \triangleleft b) \tilde{\triangleleft} b = a = (a \tilde{\triangleleft} b) \triangleleft b$	(Inv)
RI	$\leftrightarrow$	$a \triangleleft a = a$	(Idem)

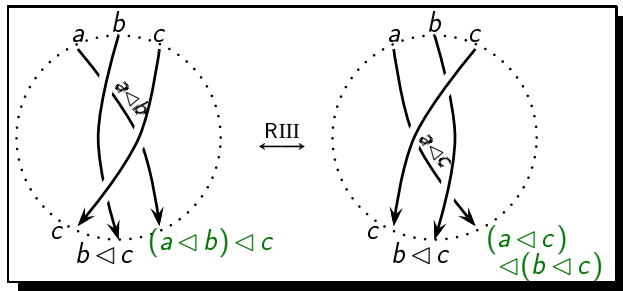
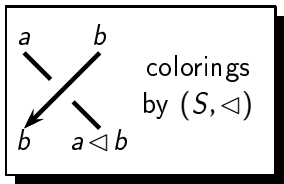
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Quandle
 (1982 D. Joyce,
 S. Matveev)

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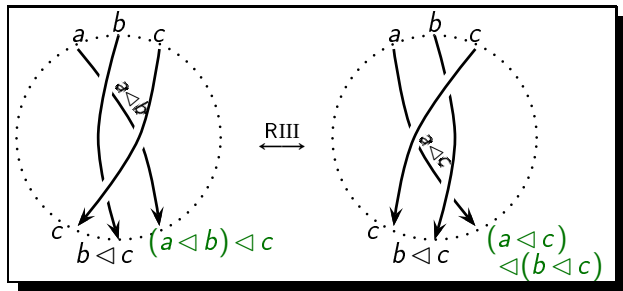
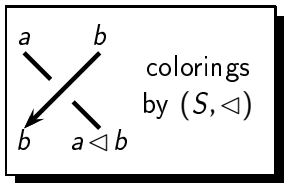


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knot invariants $\xleftrightarrow{\text{colorings}}$ quandle

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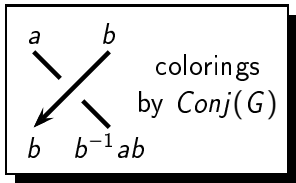
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Conjugation quandles

Group $G \rightsquigarrow$ quandle $\text{Conj}(G)$:
 $(G, g \triangleleft h = h^{-1}gh, g \widetilde{\triangleleft} h = hgh^{-1}).$

Quandles as an algebraization of knots



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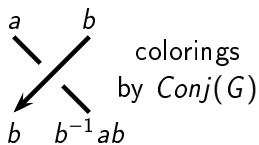
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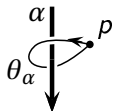
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Quandles as an algebraization of knots



Wirtinger
presentation:



colorings by $\text{Conj}(G)$

$\updownarrow$
 $\text{Rep}(\pi_1(\mathbb{R}^3 \setminus K), G)$

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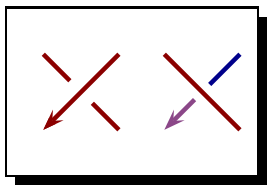
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Quandles as an algebraization of knots



colorings by R_3
 $\updownarrow$
 Fox colorings

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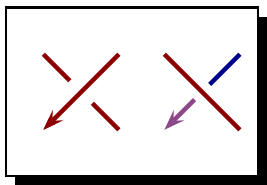
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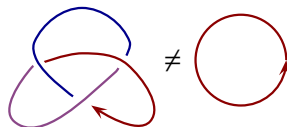
Dihedral quandles

$$R_3 = (\mathbb{Z}_3, a \triangleleft b = 2b - a).$$

Quandles as an algebraization of knots



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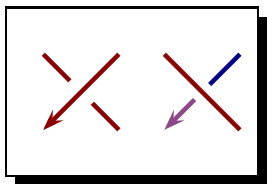
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Dihedral quandles

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Quandles as an algebraization of knots



colorings by R_n
 $\updownarrow$
 Fox n -colorings

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Digression: The fundamental quandle of a knot

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Knots $\longrightarrow$ Quandles

$K \longmapsto Q(K)$: generators $\leftrightarrow$ arcs of D_K
relations $\leftrightarrow$ crossings of D_K

$$c = a \triangleleft b$$



⚠ Does not depend on the choice of a diagram D_K of K .

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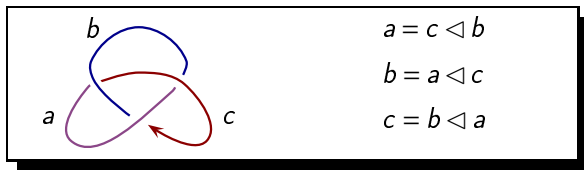
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Conclusion: The fundamental quandle is a universal weak knot invariant.

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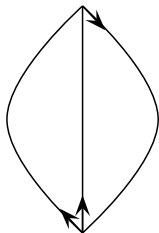
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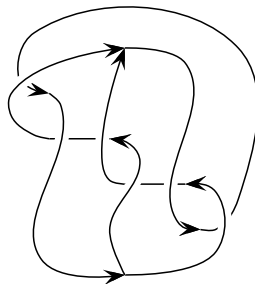
Consequence: Quandle invariants are very powerful.

Knotted 3-valent graphs

Standard and Kinoshita-Terasaka Θ -curves:



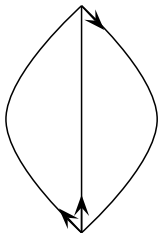
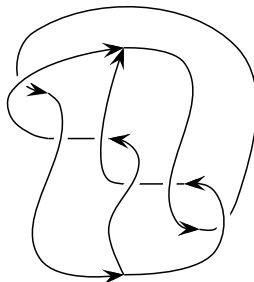
Θ_{st}



Θ_{KT}

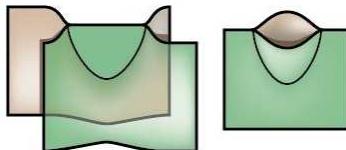
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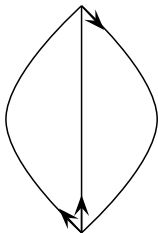
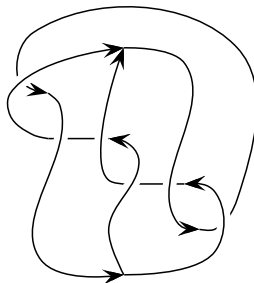
Importance:

✿ foams;



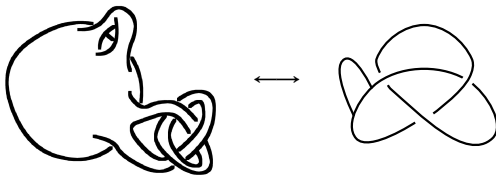
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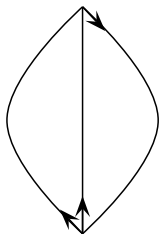
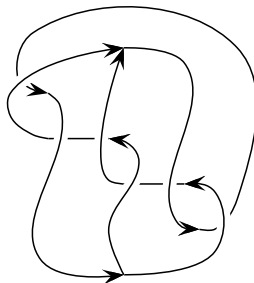
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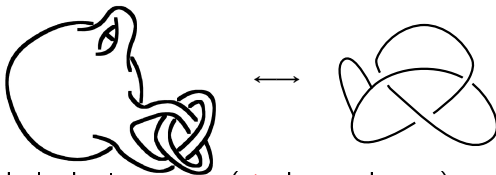
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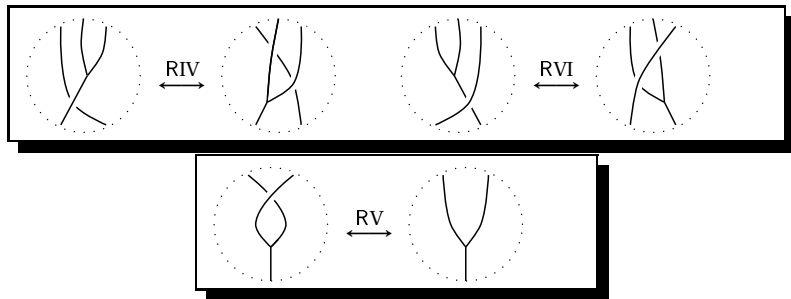
Importance:

- ✿ foams;
- ✿ handlebody-knots;
- ✿ form a finitely presented algebraic system ($\triangleleft$ knots do not).



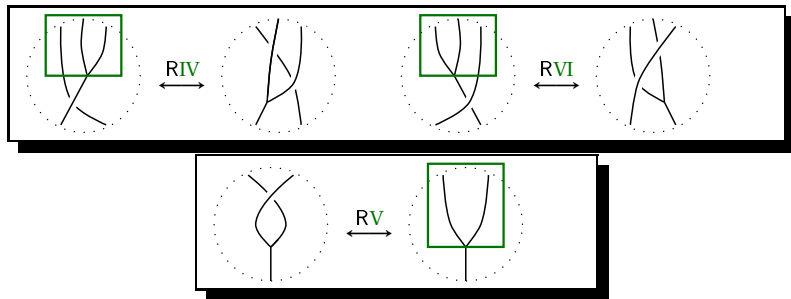
Reidemeister moves for 3-graphs

1989 L.H. Kauffman, S. Yamada, D.N. Yetter:



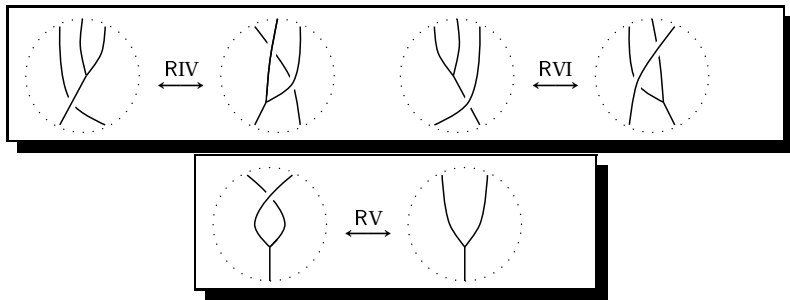
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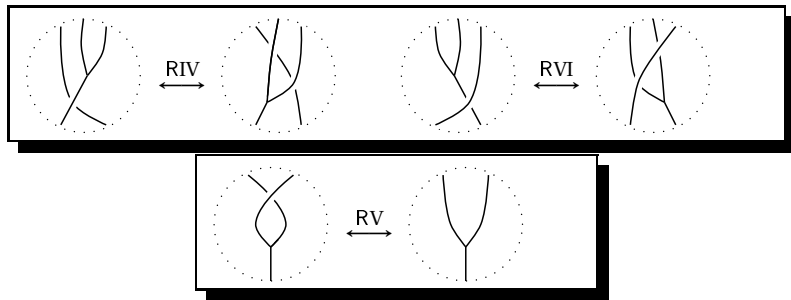
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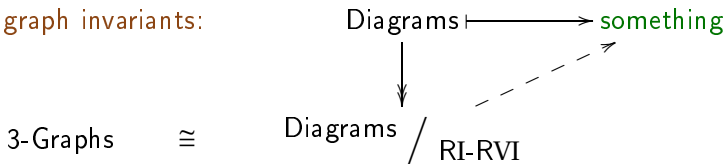
$$\text{3-Graphs} \cong \text{Diagrams} / \text{RI-RVI}$$

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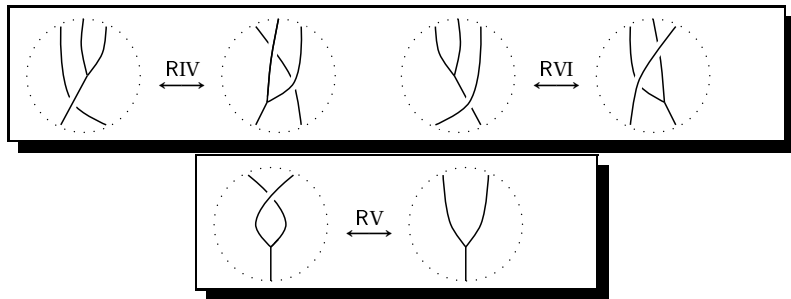


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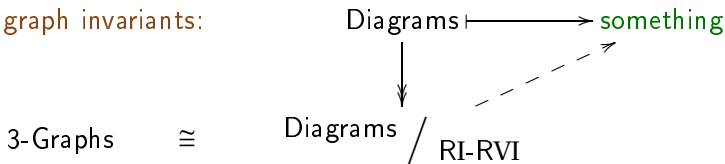


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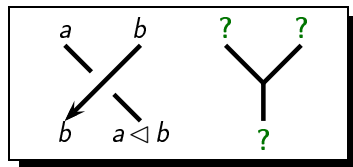
Combinatorial graph invariants:



Question: Can “something” be related to quandles?

Quandle colorings for 3-graphs

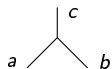
A more precise question: How to extend quandle colorings to 3-graphs?



Quandle colorings for 3-graphs

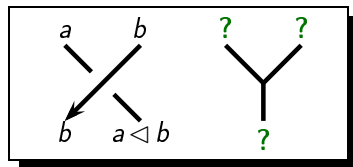
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Answers:



G is a group;

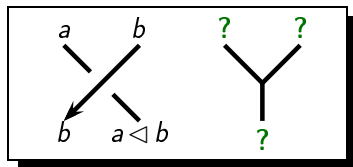
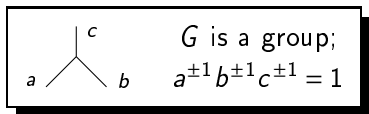
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Quandle colorings for 3-graphs

A more precise question: How to extend quandle colorings to 3-graphs?

Answers:



Generalizations:

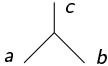
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Quandle colorings for 3-graphs

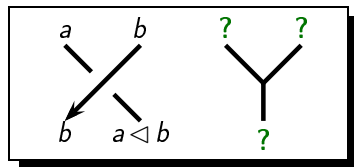
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 $\forall x, ((x \triangleleft^{\pm} a) \triangleleft^{\pm} b) \triangleleft^{\pm} c = x$

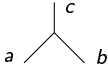
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Quandle colorings for 3-graphs

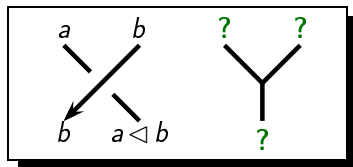
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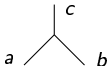
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And some more...

Orientation

Poleless 3-graphs: only *zip*  and *unzip*  vertices.

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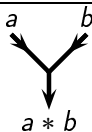
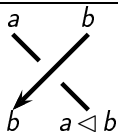
Side question: $\#$ { poleless orientations of Γ } ?

Theorem (A. Ishii): All poleless orientations of a given 3-graph are connected by single-edge orientation switches:



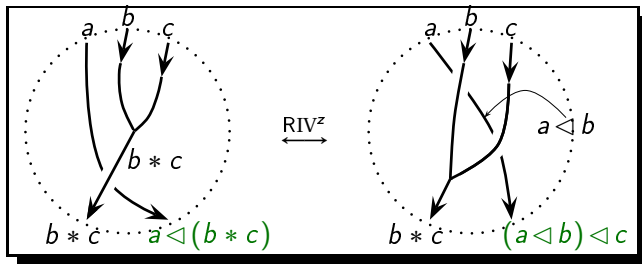
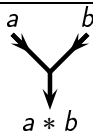
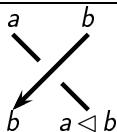
Qualgebras as an algebraization of 3-graphs

colorings
by $(S, \triangleleft, *)$

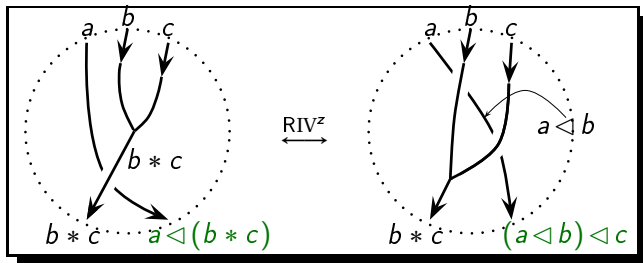
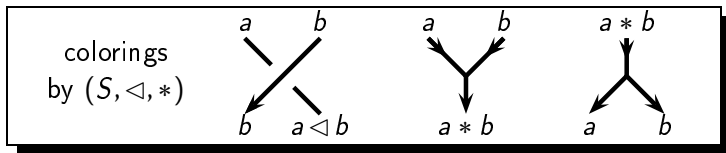


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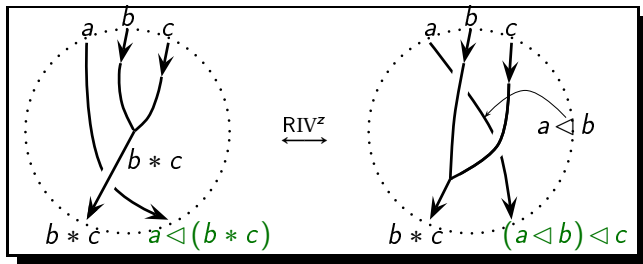
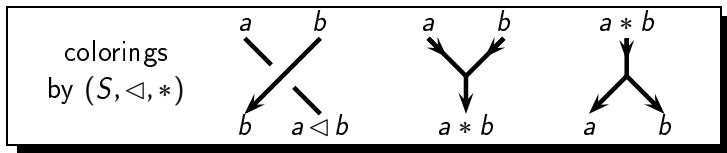


Qualgebras as an algebraization of 3-graphs



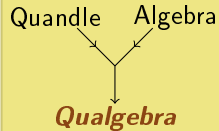
RIV	$\leftrightarrow$	$a \triangleleft (b * c) = (a \triangleleft b) \triangleleft c$
RVI	$\leftrightarrow$	$(a * b) \triangleleft c = (a * c) \triangleleft (b * c)$
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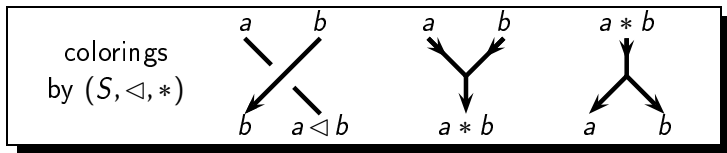


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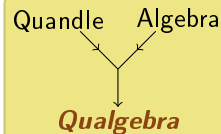


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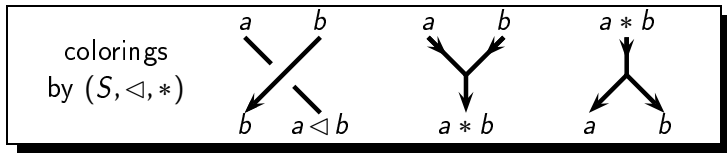
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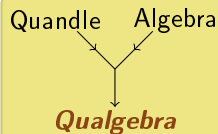
3-graph invariants $\xleftrightarrow{\text{colorings}}$ qualgebra

Qualgebras as an algebraization of 3-graphs



$$\begin{array}{ll}
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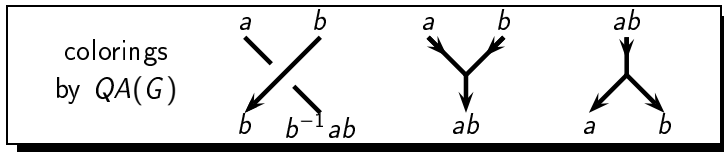


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Group qualgebras

Group $G \rightsquigarrow$ qualgebra $QA(G) = (G, g \triangleleft h = h^{-1}gh, g * h = gh)$.

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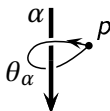


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Wirtinger
presentation:

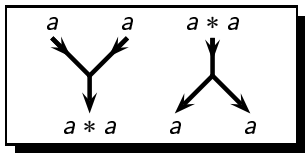


colorings by $QA(G)$

$\updownarrow$
 $Rep(\pi_1(\mathbb{R}^3 \setminus \Gamma), G)$

Computation example

Isosceles colorings:



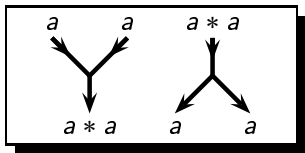
Linguistic remark:



Isosceles
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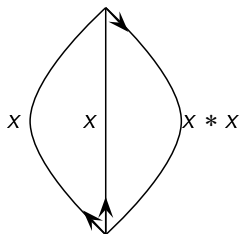
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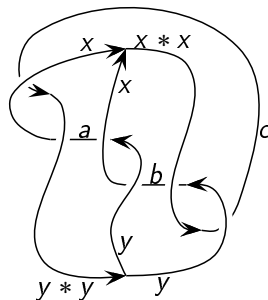
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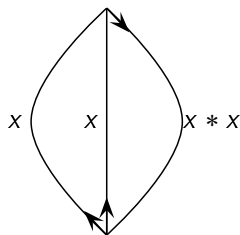
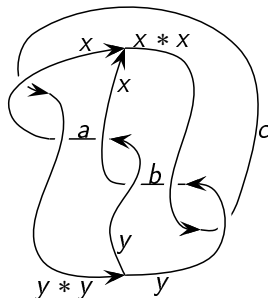


Θ_{st}



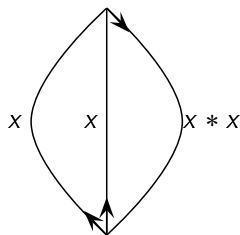
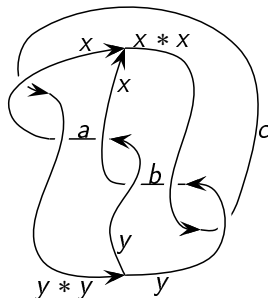
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Computation example


 Θ_{St}

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$$(\star) \begin{cases} a &= x \triangleleft (y * y) = y \triangleleft x, \\ b &= x \widetilde{\triangleleft} y = y \widetilde{\triangleleft} (x * x), \\ c &= (y * y) \triangleleft x = (x * x) \widetilde{\triangleleft} y. \end{cases}$$

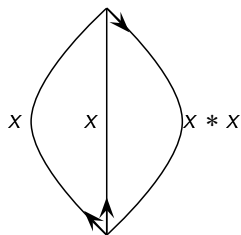
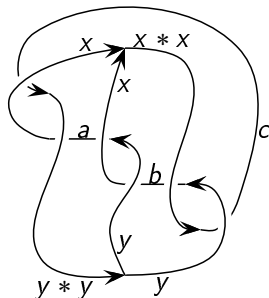
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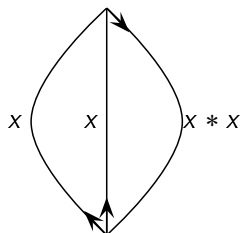
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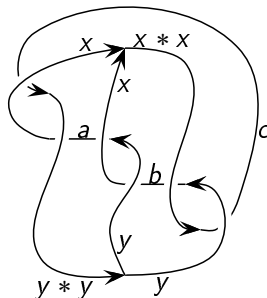
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Computation example



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$$\#Col_{S_4}^{iso}(\Theta_{St}) = \#S_4$$

$$\#Col_{S_4}^{iso}(\Theta_{KT}) \geq \#S_4 + 1$$

Part 2:

*How an Algebraist
Would Invent Qualgebras*

Group qualgebras

Example 1

Group $G \rightsquigarrow$ **conjugation quandle** $\text{Conj}(G) = (G, g \triangleleft h = h^{-1}gh)$.

1982 D. Joyce: Quandle axioms $\iff$ all properties of conjugation.

Group qualgebras

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Example: In any $QA(G)$,

$$g \triangleleft (h \triangleleft k) = ((g \widetilde{\triangleleft} k) \triangleleft h) \triangleleft k$$

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$\Rightarrow$ the equation follows from quandle axioms:

$$g \triangleleft (h \triangleleft k) \stackrel{(Inv)}{=} ((g \widetilde{\triangleleft} k) \triangleleft k) \triangleleft (h \triangleleft k) \stackrel{(SD)}{=} ((g \widetilde{\triangleleft} k) \triangleleft h) \triangleleft k.$$

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abstract level	quandle axioms	
topology	moves RI-RIII	
groups	conjugation	

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 Qualgebra axioms $\nRightarrow$ all relations btwn conjugation & multiplication.

Example: $(b \triangleleft a) * (a \triangleleft b) = ((a \widetilde{\triangleleft} b) \triangleleft a) * b$

in $QA(G)$: $a^{-1}bab^{-1}ab = a^{-1}bab^{-1}ab$

in free qualgebras: false

Other qualgebra examples

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Group G & $X \subset G \rightsquigarrow$ the sub-qualgebra of $QA(G)$ generated by X .

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$\rightsquigarrow$ Abstract graph invariants.

An infinite family of exotic examples

Example S_n

Set X & commutative operation $\star$ & zero element 0 ($0 \star x = x \star 0 = 0$)

An infinite family of exotic examples

Example S_n

Set X & commutative operation $\star$ & zero element 0 ($0 \star x = x \star 0 = 0$) $\leadsto$

$$Q_{X,n} = \{((x_1, \dots, x_n), g) \in X^{\times n} \times S_n \mid x_i = x_j = 0 \text{ whenever } g(i) = j \neq i\},$$

$$(\bar{x}, g) \triangleleft (\bar{y}, h) = (\bar{x} \cdot h, h^{-1}gh),$$

$$(\bar{x}, g) * (\bar{y}, h) = (\bar{x} \star \bar{y}, gh)$$

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$\leadsto$ two 5-element qualgebras.

An infinite family of exotic examples

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E.g., $X_2 = \{0, a\}$, $a \star_1 a = 0$ or $a \star_2 a = a$, $n = 2$, $S_2 = \{\text{Id}, \tau\}$,

$$Q_{X_2,2} = \{((x_1, x_2), \text{Id}) \mid x_1, x_2 \in X_2\} \coprod \{((0, 0), \tau)\},$$

$$\# Q_{X_2,2} = 5$$

$\leadsto$ two 5-element qualgebras.

✿ Commutative.

✿ Associative.

✿ Not cancellative $\implies$ do not come from groups:

$$((0, 0), \tau) *_i ((x_1, x_2), \text{Id}) = ((0, 0), \tau), \quad i = 1, 2.$$

Towards a classification

Example 3

Size ≤ 3 : only trivial qualgebras.

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Non-trivial qualgebra structures on $Q = \{p, q, r, s\}$:

put $\overline{p} = q, \overline{q} = p, \overline{r} = r, \overline{s} = s$;

$x \triangleleft r = \overline{x}$, $x \triangleleft y = x$ for other y ;

$*$ is commutative, $\overline{x} * \overline{y} = \overline{x * y}$,

$r * x = r$ for $x \neq r$,

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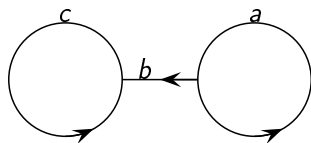
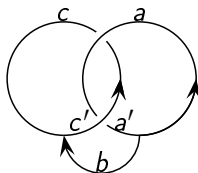
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Question: Continue the classification.

Computation example

Standard and Hopf cuff graphs:


 C_{st}


$$a' = ((a \tilde{\triangleleft} c) \triangleleft a)$$

$$c' = c \triangleleft a$$

 C_H

$$\#Col_Q(C_{st}) = \#\{(a, b, c) \in Q \mid b * a = a, b * c = c\} = 18,$$

$$\#Col_Q(C_H) = \#\{(a, b, c) \in Q \mid b * a = a \triangleleft c, b * c = c \triangleleft a\} = 14.$$

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Group $G \rightsquigarrow QA^*(G) = (G, g \triangleleft h = h^{-1}gh, g * h = gh, \rho(h) = h^{-1}).$

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- $\rightsquigarrow *$ is a **Latin square** (= pseudo-sudoku).

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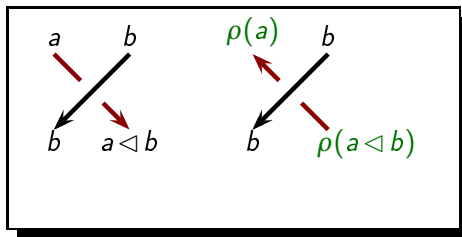
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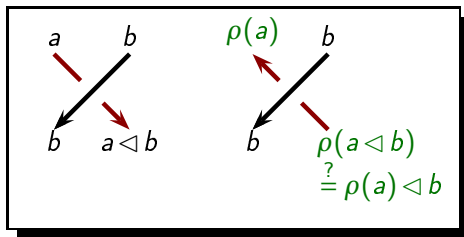
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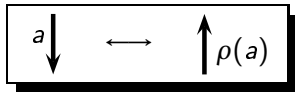
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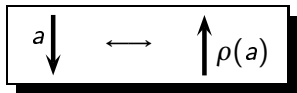
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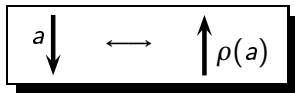
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abstract level	good involution axioms for qualgebras
topology	<u>un</u> oriented 3-graphs
groups	conjugation- and multiplication-inversion interactions

Symmetric qualgebras: examples

Example 0

Symmetric trivial qualgebras $\longleftrightarrow$ Latin squares which

- ✿ are symmetric w.r.t. the main diagonal, and
- ✿ contain a row $\sigma \in \text{Bij}(S)$ iff contain a row σ^{-1} .

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Example 3

$S = \{x, y, z\}$, $a \triangleleft b = a$, $*$ is commutative.

$*$	x	y	z
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y	y	z	x
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$QA(\mathbb{Z}_3)$

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✿ Trivial qualgebras:

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$\leftarrow QA(\mathbb{Z}_4)$

$QA(\mathbb{Z}_2 \times \mathbb{Z}_2) \rightarrow$

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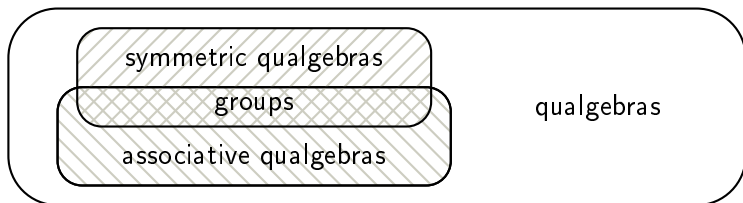
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Alexander quandle $(M, a \triangleleft b = \alpha a + (1 - \alpha)b)$, $M \in {}_R\text{Mod}$, $\alpha \in R^*$.

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Uniqueness

$QA(S_n)$ is the unique qualgebraization of $\text{Conj}(S_n)$ ($\iff T$ is injective).

From quandles to qualgebras: examples

Size 3

- 1 Trivial quandle: 129 qualgebraizations.

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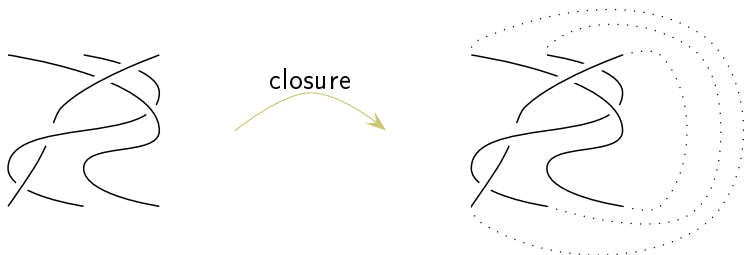
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Part 3:

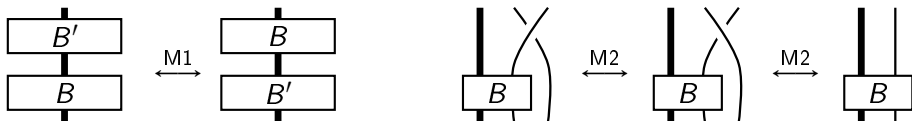
*Variations of
Qualgebra Ideas*

Alexander-Markov theorem



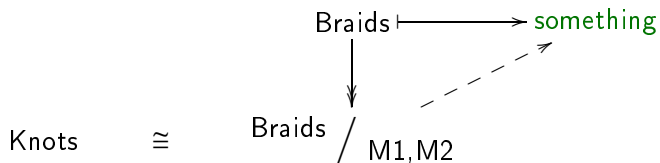
Theorem (1923 Alexander; 1935 Markov)

- ✿ Surjectivity.
- ✿ Kernel: moves M1 and M2.



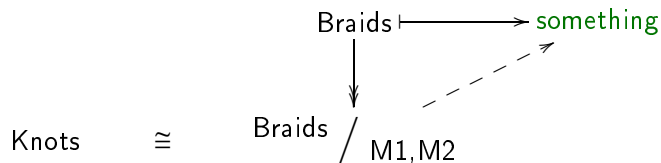
Braid and knot invariants

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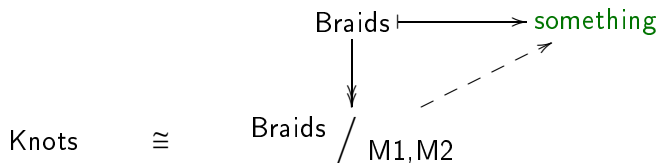


1925 E. Artin:

$$\text{Braids} \cong \text{Diagrams} / \text{RII-RIII}$$

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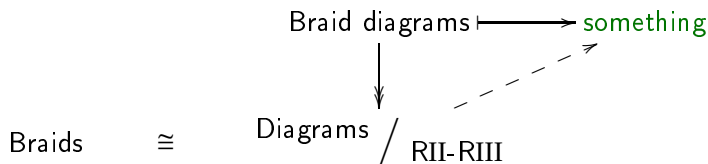
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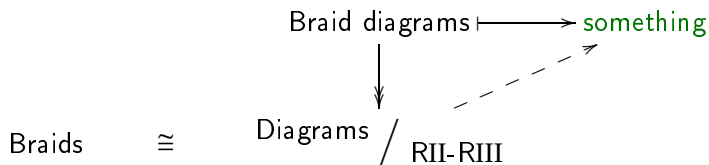
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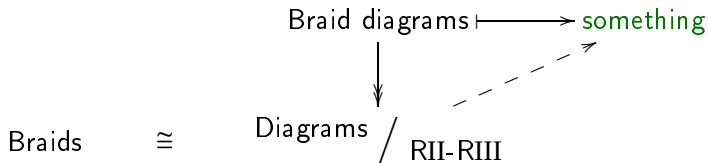


Many knot invariants adapt to the braid case, with possible **enhancements**.

Example: quandle invariants.

Braid and knot invariants

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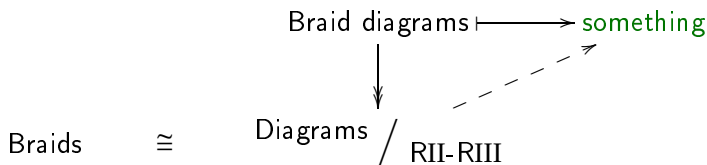
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✿ Weaker structure: **rack** (= quandle without $a \triangleleft a = a$) $\iff$ no RI.

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✿ **Operator invariants** instead of counting invariants:

braids with n strands $\longrightarrow \text{Aut}(S^{\times n})$,

$\beta \longmapsto$ (colors of upper arcs
 $\mapsto$ colors of lower arcs).

Branched Alexander-Markov theorem

Question: : A closure procedure for 3-graphs?

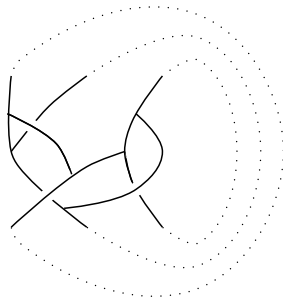
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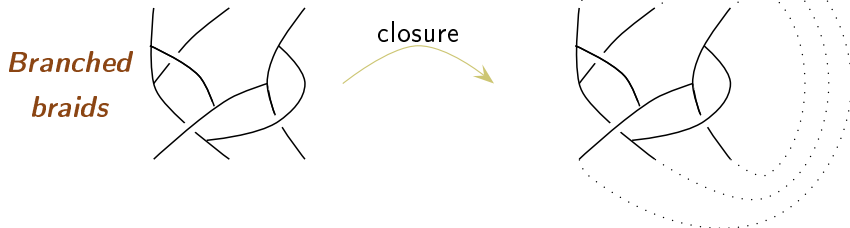


closure



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Theorem (2010 K. Kanno - K. Taniyama; 2014 S. Kamada - L.)

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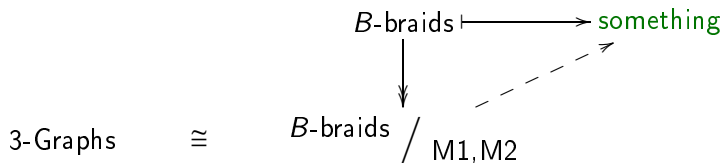


Generalizations

- ✿ Graph-braids (vertices of arbitrary valence).
- ✿ Virtual and welded versions.

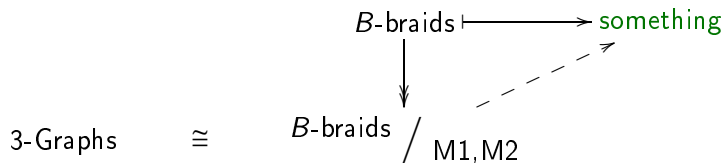
Branched braid and 3-graph invariants

3-Graph invariants out of B -braid invariants:

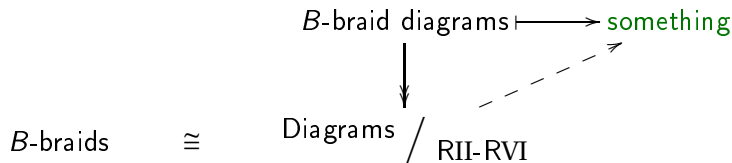


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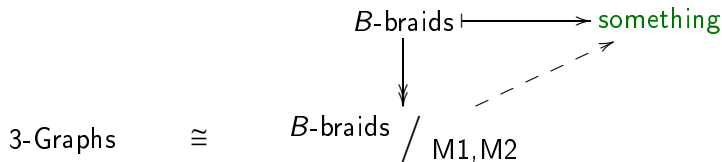


Combinatorial B -braid invariants:

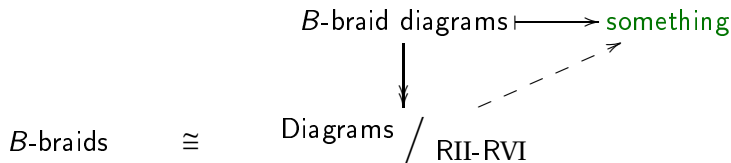


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B -braid invariants $\xrightarrow{\text{colorings}}$ **weak qualgebra**

(omit $a \triangleleft a = a$)

Getting more out of quaugebra colorings

diagrams:	D	$\xrightarrow[\rightsquigarrow]{\text{R-move}}$	D'
colorings:	$\mathcal{C}$	$\rightsquigarrow$	$\mathcal{C}'$
coloring sets:	$\text{Col}_S(D)$	$\xleftrightarrow{\text{bij}}$	$\text{Col}_S(D')$

Getting more out of quaugebra colorings

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Counting invariants:

$$\#\text{Col}_S(D) = \#\text{Col}_S(D').$$

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Getting more out of quandle colorings

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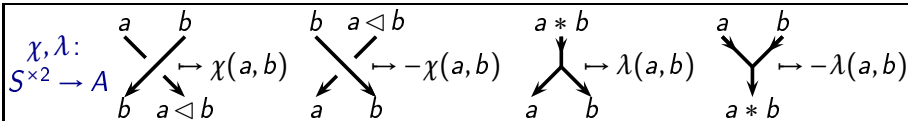
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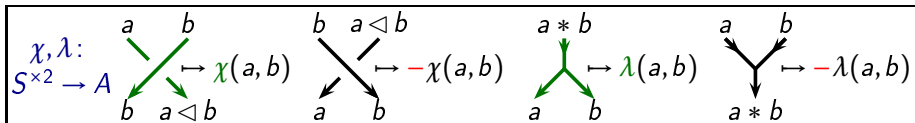
Inspiration: *Quandle cocycle invariants* of knots (1999 Carter-Jelsovsky-Kamada-Langford-Saito).

Qualgebra cocycle invariants for 3-graphs



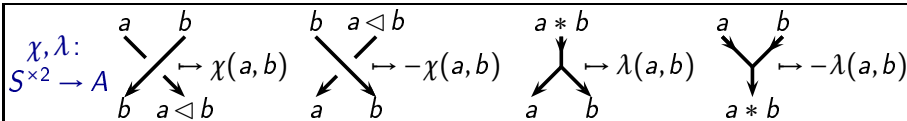
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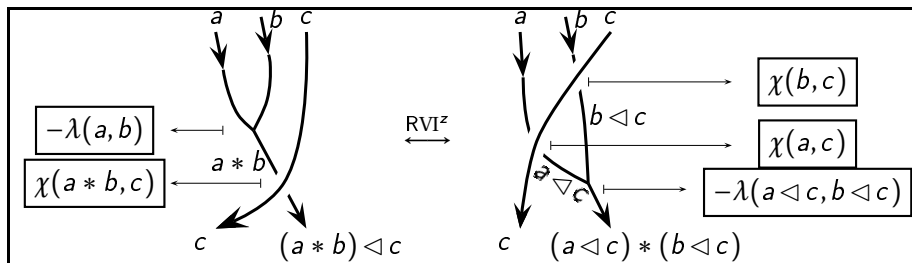


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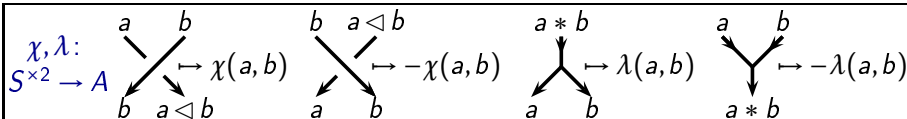
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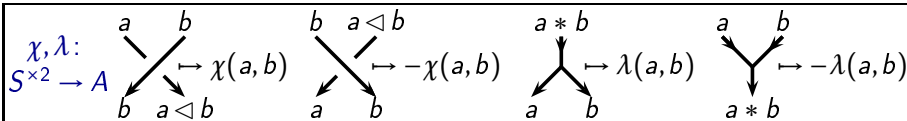
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**Qualgebra
2-cocycle**

RI-RIII are automatic.

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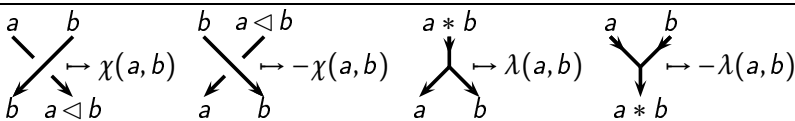
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Qualgebra cocycle invariants $\supseteq$ qualgebra counting invariants

($\rightsquigarrow \chi = \lambda = 0$).

Towards qualgebra cohomology

Qualgebra 2-coboundaries ($\subseteq$ qualgebra 2-cocycles):

$$\begin{aligned} \phi: S \rightarrow A \quad \rightsquigarrow \quad \chi(a, b) = \phi(a \triangleleft b) - \phi(a), & \quad \rightsquigarrow \quad \text{trivial} \\ \lambda(a, b) = \phi(a) + \phi(b) - \phi(a * b) & \quad \text{graph invariants} \end{aligned}$$

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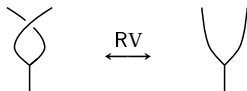
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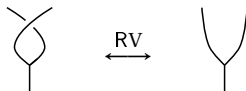
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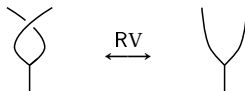
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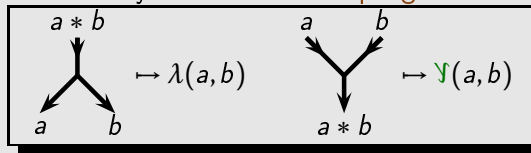
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Enhancements

✿ **Region colorings** and shadow cocycle invariants $\leadsto$ **qualgebra 3-cocycles**.

✿ Distinguish zip- and unzip-vertices:



Quasialgebra cocycles: example

Example 4

$$A = \mathbb{Z}, \quad Q = \{p, q, r, s\}$$

$$\bar{p} = q, \bar{q} = p, \bar{r} = r, \bar{s} = s;$$

$$x \triangleleft r = \bar{x}, \quad x \triangleleft y = x \text{ for other } y;$$

$$* \text{ is commutative,} \quad \bar{x} * \bar{y} = \overline{x * y},$$

$$r * x = r \text{ for } x \neq r,$$

$$r * r = s * s = p * q = s, \quad p * p, p * s \in \{p, q, s\}$$

$$\ast Z^2(Q) \cong \mathbb{Z}^8$$

$$\ast B^2(Q) \cong \mathbb{Z}^4$$

$$\ast H^2(Q) \cong \mathbb{Z}_2 \oplus \mathbb{Z}^4$$



どうもありがとう