

MA3466 Tutorial Sheet 7, q1-2 solutions<sup>1</sup>

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1. (C& T 5.2) Let  $p(S_i) = p_i$  for some set of outcomes  $\{S_1, S_2, \dots, S_n\}$ . The  $S_i$ 's are uniquely encoded into strings from a  $D$ -symbol alphabet in a uniquely decodable manner. If  $n = 6$  and the code word lengths are  $(l_1, l_2, \dots, l_n) = (1, 1, 2, 3, 2, 3)$  find a good lower bound on  $D$ .

*Solution:* Well this question is easy to do by checking the first few cases.  $D = 2$  clearly doesn't work, since  $l(S_1) = l(S_2) = 1$  then we need to have  $c(S_1) = 0$  and  $c(S_2) = 1$  or visa versa, and then the length two code word for  $S_3$ , which needs to be 00, 01, 10 or 11, is ambiguous, since it could also be decoded as a pair from  $S_1$  and  $S_2$ . However,  $D = 3$  does work, for example

	$S_1$	$S_2$	$S_3$	$S_4$	$S_5$	$S_6$
	0	1	20	220	21	221

is a prefix code with the correct lengths.

2. (C& T 5.4) Slackness in the Kraft inequality. An instantaneous code has word lengths  $l_1$  to  $l_m$  satisfying the strict inequality

$$\sum_{i=1}^m D^{-l_i} < 1 \tag{1}$$

Show there are arbitrarily long sequences of code symbols in  $\mathcal{D}^*$  which cannot be decoded into sequences of codewords: that is, not all sequences of symbols in  $D$  form a sentence.

*Solution:* So the trick here is that, if

$$\sum_{i=1}^m D^{-l_i} < 1 \tag{2}$$

then, with  $l_x$  with maximum length,

$$\sum_{i=1}^m D^{l_x - l_i} < D^{l_x} \tag{3}$$

and, we know, that  $D^{l_x}$  is the number of nodes at level  $l_x$ , which the sum on the right hand side gives the number of nodes at that level which are codewords or their

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descendants. Thus, the strict inequality means that there are nodes at level  $l_x$  which are not codewords and which are not the descendants of codewords. Taking one of these nodes: the node will correspond to a codeword and could have been added to the existing prefix code to extend it to a further symbol, this code would still be uniquely decodable. This means that strings in the codeword corresponding to this node do not correspond to any of the existing symbols.