## MA3466 Tutorial Sheet 3, outline solutions<sup>1</sup>

## $22\ {\rm March}\ 2010$

- 1. (C&T 2.2) Entropy of functions. Let X be a random variable taking on a finite number of values. What is the general inequality relating H(X) and H(Y) if
  - (a)  $Y = 2^X$
  - (b)  $Y = \cos X$

Solution: The key point here is that if Y = g(X) then X determines Y but the converse may or may not be true. This means that H(Y|X) = 0 since Y is not random is the outcome of X is known. However H(X|Y) may not be zero if the function is not invertible.

In the first example it is true because the function is one-to-one. First, because  $Y=2^{X}$ 

$$p(Y = 2^x | X = x) = 1 \tag{1}$$

and so H(X, Y) = H(X) + H(Y|X) = H(X); H(Y|X) = 0 because

$$H(Y|X) = \sum p(x)H(Y|X=x)$$
<sup>(2)</sup>

and all the terms in the sum in

$$H(Y|X = x) = \sum p(y|x) \log p(y|x)$$
(3)

are zero, either because  $y = 2^x$  so that p(y|x) = 1 and its log is zero, or  $y \neq 2^x$  and  $x \log x$  goes to zero as x goes to zero. The converse is also true,  $X = \log X$  and so

$$H(X, Y) = H(Y) + H(X|Y) = H(Y)$$
 (4)

and so H(X) = H(Y). The situation is different for  $Y = \cos X$ , this is not, in general, an invertable function. Hence, it is still true that H(Y|X) = 0 because X still determines Y; however H(X|Y) may not be zero, there may be y such that the set  $\cos^{-1}y = \{x \in \mathcal{X} | \cos x = y\}$  may have more than one element and so knowing Y = y tells you  $x \in \cos^{-1}y$ , but it doesn't tell you what X is. Hence

$$H(X) = H(X, Y) = H(Y) + H(X|Y) \ge H(Y)$$
 (5)

where we know  $H(X|Y) \ge 0$  because entropy is always positive.

To illustrate this further, lets consider two examples, first, if  $\mathcal{X} = \{0, \pi\}$  then the function is invertible,  $\mathcal{Y} = \{1, -1\}$  and if Y = 1, X = 0, if Y = -1,  $X = \pi$ . Here the inequality will be sharp. On the other hand, say  $\mathcal{X} = \{0, \pi, 2\pi\}$  and

$$p_X(0) = p_X(\pi) = p_X(2\pi) = 1/3 \tag{6}$$

now  $\mathcal{Y} = \{1, -1\}$  with  $p_Y(1) = 2/3$  and  $p_Y(-1) = 1/3$ . Hence

$$H(X) = \log 3 H(Y) = \log 3 - \frac{2}{3}$$
(7)

and H(Y) < H(X); the point being that  $H(X|Y) \neq 0$ , if Y = 1, X could be zero or  $2\pi$  with equal probability so H(X|Y = 1) = 1 and

$$H(X|Y) = \frac{2}{3}H(X|Y=1) + \frac{1}{3}H(X|Y=-1) = \frac{2}{3}$$
(8)

2. (C&T 2.4) Entropy of functions of a random variable. Let X be a discrete random variable. Show that the entropy of a function of X is less than or equal to the entropy of X by justifying the following steps

$$\begin{aligned} H(X,g(X)) &= H(X) + H(g(X)|X) \\ &= H(X), \\ H(X,g(X)) &= H(g(X)) + H(X|g(X)) \ge H(g(X)) \end{aligned}$$
(9)

and hence  $H(g(X)) \leq H(X)$ .

Solution: The key point here is that if Y = g(X) then X determines Y but the converse may or may not be true. In the first example it is true because the function is one-to-one. First, because  $Y = 2^X$ 

$$p(Y = 2^x | X = x) = 1 \tag{10}$$

and so H(X,Y) = H(X) + H(Y|X) = H(X); H(Y|X) = 0 because

$$H(Y|X) = \sum p(x)H(Y|X=x)$$
(11)

and all the terms in the sum in

$$H(Y|X = x) = \sum p(y|x) \log p(y|x)$$
(12)

are zero, either because  $y = 2^x$  so that p(y|x) = 1 and its log is zero, or  $y \neq 2^x$  and  $x \log x$  goes to zero as x goes to zero. The converse is also true,  $X = \log X$  and so

$$H(X,Y) = H(Y) + H(X|Y) = H(Y)$$
 (13)

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and so H(X) = H(Y). The situation is different for  $Y = \cos X$ , this is not, in general, an invertable function. Hence, it is still true that H(Y|X) = 0 because X still determines Y; however H(X|Y) may not be zero, there may be y such that the set  $\cos^{-1}y = \{x \in \mathcal{X} | \cos x = y\}$  may have more than one element and so knowing Y = y tells you  $x \in \cos^{-1}y$ , but it doesn't tell you what X is. Hence

$$H(X) = H(X,Y) = H(Y) + H(X|Y) \ge H(Y)$$
 (14)

where we know  $H(X|Y) \ge 0$  because entropy is always positive.

To illustrate this further, lets consider two examples, first, if  $\mathcal{X} = \{0, \pi\}$  then the function is invertible,  $\mathcal{Y} = \{1, -1\}$  and if Y = 1, X = 0, if Y = -1,  $X = \pi$ . Here the inequality will be sharp. On the other hand, say  $\mathcal{X} = \{0, \pi, 2\pi\}$  and

$$p_X(0) = p_X(\pi) = p_X(2\pi) = 1/3 \tag{15}$$

now  $\mathcal{Y} = \{1, -1\}$  with  $p_Y(1) = 2/3$  and  $p_Y(-1) = 1/3$ . Hence

$$H(X) = \log 3$$
  
 $H(Y) = \log 3 - \frac{2}{3}$ 
(16)

and H(Y) < H(X); the point being that  $H(X|Y) \neq 0$ , if Y = 1, X could be zero or  $2\pi$  with equal probability so H(X|Y = 1) = 1 and

$$H(X|Y) = \frac{2}{3}H(X|Y=1) + \frac{1}{3}H(X|Y=-1) = \frac{2}{3}$$
(17)

3. (C&T 2.8) Drawing with and without replacement. An urn contains r red, w white and b black balls. Which has higher entropy, drawing  $k \ge 2$  balls from the urn with replacement or without replacement?

Solution: So the answer to this question relies on the fact that the probability distribution for the n drawing is the same irrespive of whether there is replacement or not. Lets use X to denote drawing from an urn with r red balls, w white balls and b black balls, so, with n = b + r + w

$$p_X(c_r) = \frac{r}{n}$$

$$p_X(c_w) = \frac{w}{n}$$

$$p_X(c_b) = \frac{b}{n}$$
(18)

whre  $c_r$  is red and so on. Now, if  $X_i$  is the *i*th drawing with replacement, then clearly the  $X_i$  are independent and  $p_{X_i}(x) = p_X(x)$  for  $x \in \mathcal{X} = \{c_r, c_b, c_w\}$ .

Now, let  $Y_i$  be the *i*th drawing with replacement: although the  $Y_i$  are not independent  $p_{Y_i}(x) = p_X(x)$  for  $x \in \mathcal{X}$ . To see this, note  $Y_1 = X$  and assume it is true for  $Y_i$  and consider  $Y_{i+1}$ :

$$p_{Y_{i+1}}(c_r) = p_{(Y_{i+1},Y_i)}(c_r,c_r) + p_{(Y_{i+1},Y_i)}(c_r,c_w) + p_{(Y_{i+1},Y_i)}(c_r,c_b)$$

$$= p_{Y_{i+1}|Y_i}(c_r|c_r)p_{Y_i}(c_r) + p_{Y_{i+1}|Y_i}(c_r|c_w)p_{Y_i}(c_w) + p_{Y_{i+1}|Y_i}(c_r|c_b)p_{Y_i}(c_b) = \frac{r-1}{n-1}\frac{r}{n} + \frac{r}{n-1}\frac{w}{n} + \frac{r}{n-1}\frac{b}{n} = \frac{r}{n} = p_X(c_r)$$
(19)

This means, using the chain rule and the conditioning theorem

$$\begin{aligned} H(Y_1, Y_2, \dots, Y_n) &= H(Y_1) + H(Y_2|Y_1) + H(Y_3|Y_2, Y_1) + \dots + H(Y_n|Y_{n-1}, \dots, Y_1) \\ &\leq \sum H(Y_i) = nH(X) = H(X_1, X_2, \dots, X_n) \end{aligned}$$

with equality if and only if the  $Y_i$  were independent which they aren't, hence

$$H(Y_1, Y_2, \dots, Y_n) < H(X_1, X_2, \dots, X_n)$$
 (21)

- 4. (C&T 2.14) Enropy of a sum. Let X and Y be random variables that take on values  $x_1, x_2, \ldots, x_r$  and  $y_1, y_2, \ldots, y_s$  respectively. Let Z = X + Y.
  - (a) Show that H(Z|X) = H(Y|X). Argue that if X and Y are independent then  $H(Y) \leq H(Z)$  and  $H(X) \leq H(Z)$ . Thus the addition of independent random variables add uncertainy.
  - (b) Give an example of random variables for which H(X) > H(Z) and H(Y) > H(Z).
  - (c) Under what conditions does H(Z) = H(X) + H(Y).

Solution: So, given X, Y determines Z and visa versa, so H(Z|X) = H(Y|X). Now, we know that

$$H(Y|X) = H(Z|X) \le H(Z) \tag{22}$$

but, if X and Y are independent, H(Y|X) = H(Y), so  $H(Y) \le H(Z)$ ;  $H(X) \le H(Z)$ follow by a similar argument. Thus, if we want H(X) > H(Z), we need X and Y dependent. In fact, we want X and Y to be dependent in such a way that adding them gives something less uncertain; as an example, let Y = -X so Z = 0 always and so, H(Z) = 0 and is less than H(X) = H(Y) for any non-trivial choice of X. Finally,

$$H(X+Y) = H(X,Y) \tag{23}$$

if the addition is invertible, that is, if there are unique X and Y for any X + Y; this would happen, for example, if  $\mathcal{X} = \{1, 2\}$  and  $\mathcal{Y} = \{1, 3\}$  since the possible values of the sum are 2, 3, 4 and 5 and each correponds to a different choice of X and Y; however, if  $\mathcal{X} = \{1, 2\}$  and  $\mathcal{Y} = \{1, 2\}$  then X = 1, Y = 2 and X = 2, Y = 1 both give X + Y = 3. Now

$$H(X,Y) = H(X) + H(Y|X)$$
(24)

and H(Y|X) + H(Y) if X and Y are independent.