MA3466 Tutorial Sheet 1, outline solutions¹

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1. For two random variables with numerical outcomes, find p(x, y) so that there is zero correllation

$$x'y'\rangle = 0\tag{1}$$

but X and Y aren't independent. x' = x - EX and y' = y - EY.

Solution:

I tried this first for the case where $|\mathcal{X}| = |\mathcal{Y}| = 2$ and found in this case C = 0 implies $p_{X,Y}(x,y) = p_X(x)p_Y(y)$. In fact, a simple parameter counting might be enough here, but a full argument goes like, consider p(x, y) given by

with X across and Y down. Now, working out the marginal distributions we have $p_X(0) = a + c$, $p_X(1) = b + d$ so that $\langle X \rangle = b + d$; $p_Y(0) = a + b$ and $p_Y(1) = c + d$ so $\langle Y \rangle = c + d$. Substituting into the formula for C we get

$$C = a(-b-d)(-c-d) + b(1-b-d)(-c-d) +c(-b-d)(1-c-d) + d(1-b-d)(1-c-d) = d - (b+d)(c+d)$$
(3)

so C = 0 means d = (b + d)(c + d). Of course, since these are probabilities, we also have a + b + c + d = 1. Now, this probability table corresponds to independent distributions, p(x, y) = p(x)p(y), if the columns of the table are proportional to each other, since p(x, y = 0) = p(x)p(y = 0) and p(x, y = 1) = p(x)p(y = 1) = p(x, y = 0)p(y = 1)/p(y = 0). This implies ac = bd, substituting a = 1 - b - c - d into this equation gives d = (c + d)(b + d), precisely the condition for C = 0, zero correlation implies statistical independence for two by two probability tables.

Next, lets set $|\mathcal{X}| = 2$ and $|\mathcal{Y}| = 3$. Rather than look at the general case, consider

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It is easy to check that this already has C = 0, but the two columns are not proportional provided $ad \neq bc$, showing that there are distributions with zero correlation which are statistically independent. As a concrete example take

2. Work out the marginal distributions and the x = a conditional distribution for

	a	b
1	$\frac{1}{3}$	<u>1</u> 6
2	Ŏ	$\frac{1}{4}$
3	$\frac{1}{8}$	$\frac{1}{8}$

Solution: Calculating the marginal distribution just requires adding along the rows or columns, if we take X to be the random variable going across, so $\mathcal{X} = \{a, b\}$ and Y the random variable going down, so $\mathcal{Y} = \{1, 2, 3\}$ then the two marginal distributions are, for X

$$\begin{array}{c|cc} a & b \\ \hline 11 \\ \hline 24 \\ \hline 24 \\ \end{array}$$

and for Y

The conditional distribution, conditioned on x = a is calculated using p(y|x)p(x) = p(x, y) so we divide the p(x = a, y) column by p(x = a) = 11/24, hence

3. (C&T 2.1) A fair coin is flipped until the first head occurs. Let X denote the number of flips required.

(a) Find the entropy H(X) in bits. The following expressions may be useful:

$$\sum_{n=0}^{\infty} r^{n} = \frac{1}{1-r}$$

$$\sum_{n=0}^{\infty} nr^{n} = \frac{r}{(1-r)^{2}}$$
(6)

(b) A random variable X is drawn according to this distribution. Find an efficient sequence of yes-no questions of the form, 'Is X contained in the set S?'. Compare H(X) to the expected number of questions required to determine X.

Solution: So, here the set of possible outcomes is $\mathcal{X} = \{1, 2, 3, ...\}$ and we need to start by working out $p_X n$ the chance of throwing n flips before getting a head. To get X = n you need to throw n - 1 tails, this has probability $1/2^{n-1}$ followed by a head, which has probability 1/2, hence

$$p_X(n) = \frac{1}{2^n} \tag{7}$$

It is easy to check that

$$\sum_{n=1}^{\infty} p_X(n) = \sum_{n=1}^{\infty} \frac{1}{2^n} = \sum_{n=0}^{\infty} \frac{1}{2^n} - 1 = 1$$
(8)

Now, to calculate the entropy, just use the formula:

$$H(X) = -\sum_{n \in \mathcal{X}} p_X(x) \log p_X(x)$$

= $-\sum_{n=1}^{\infty} \frac{1}{2^n} \log 2^{-n} = \sum_{n=1}^{\infty} \frac{n}{2^n}$
= 2 (9)

It is easy to see that this is the same as the average number of questions asked in order starting with n = 1 of the form 'is the answer n?' it would take to find X.

4. (C&T 2.3) What is the minimum value of $H(p_1, p_2, \ldots, p_n) = H(\mathbf{p})$ as \mathbf{p} ranges overall possible vectors. Find the \mathbf{p} which achieve this bound. $H(p_1, p_2, \ldots, p_n)$ is a common notation for H(X) where X has n-outcomes $\{x_1, x_2, \ldots, x_n\}$ and $p_1 = p(x_1)$, $p_2 = p(x_2)$ and so on.

Solution: First, let's extremize

$$H = -\sum p_i \log p_i \tag{10}$$

subject to the contraint $\sum p_i = 1$: we will ignore the inequality constraints $p_i \ge 0$ for now. Using a Lagrange multiplier, we define

$$L = -\sum p_i \log p_i + \lambda \left(\sum p_i - 1\right) \tag{11}$$

and so, rembering that we have logs to the base two

$$\frac{\partial L}{\partial p_i} = -\log p_i - \log_2 e + \lambda \tag{12}$$

Setting $\partial L/\partial p_i = 0$ shows the p_i are all the same and hence equal to 1/n, solving for λ gives the same answer. Hence, an extreme value of H is given by $H = \log n$, this, however, is the maximum. The minima don't have $\partial L/\partial p_i = 0$, rather they lie in the corners of the inequality constraint region, by linear programming, or by examining H it is easy to see H = 0 if and only if $p_i = 1$ for some i and $p_j = 0$ for $i \neq j$; since H > 0 these are the only minima, and there are n of them.